Last time

Set N of n voters, A of m alternatives / candidates
Each voter has a strict linear order $\succ_i \in L(A)$. (= preference / ranking)
Profile $P = (\succ_1, \dots, \succ_n)$

Social Choice Function $f : P \to A$
" " Correspondence $F : P \to 2^A$
Welfare Function $F : P \to L(A)$

A SCF is strategyproof (SP) if $\forall P$, $\succ_i$, $\succ_i'$, $\succ_{-i}$,
$f(\succ_i, \succ_{-i}) \succeq_i f(\succ_i', \succ_{-i})$

Gib-Sat Theorem For $m \geq 3$, every SCF that is onto & strategyproof is a dictatorship.

strategyproof & monotone
If $|L(A)| = 2$ monotone if let $f(P) = a$, and P is a 0- improvement of P. Then $f(P) = a$.

Obvious f is dictatorial $\Rightarrow f$ is monotone

Notation : For $m \geq 3$, every SCF that is onto & satisfies Maskin monotonicity is a dictatorship.

Single - Peaked Preferences

called the "basic order"
Given a set of alternatives A, let $\triangleright \in L(A)$ be a linear order.
A voter $\succ$ is single-peaked wrt $\triangleright$ if:
Let $\succeq_{\triangleright}$ top rated alternative for $\succ$ Then:
(i) $\forall\, x \triangleright y \triangleright \succeq_{\triangleright}$, $x \succ y$
(ii) $\forall\, y \triangleright x \triangleright \succeq_{\triangleright}$, $x \succ y$

Then $\succeq_{\triangleright}$ is the "peak". As moving away from $\succeq_{\triangleright}$ makes a candidate worse

Obvious If $\succ$ is single-peaked wrt $\triangleright$, then the last-ranked alternative in $\succ$ is either the first or the last-ranked alternative in $\triangleright$

Def : A preference profile $P = (\succ_i)_{i \in N}$ is single-peaked wrt $\triangleright$ if each voter $\succ_i$ is single-peaked wrt $\triangleright$.

Lemma : Given a profile P, $N(a,b) = |\{i : a \succ_i b\}|$
$\mu(a,b) = N(a,b) - N(b,a)$
a "beats" b if $\mu(a,b) > 0$
a is a Condorcet winner if $\forall\, b \neq a$, a beats b

Thm : P: $a \succ b \succ c$
$b \succ c \succ a$
$c \succ a \succ b$ has a cycle, hence no Condorcet winner.

But this is not a single-peaked profile, for any base order.

Recognizing Single - peaked profiles

Obvious A voter $\succ_i$ is single-peaked wrt $\triangleright$ iff $\forall\, k = 1, \dots, m$, the top k alternatives in $\succ_i$ form an interval in $\triangleright$.
(term paradigm)

CSP :
Let $A \in \{0,1\}^{m \times m}$. Then A has the "consecutive 1's property" (CSP) if $\exists$ a permutation of the columns, so that in each row, the 1's are consecutive.

Theorem (Booth & Lueker 1976) : Given $A \in \{0,1\}^{m \times m}$, we can determine in polynomial time if A has the CSP, and if yes, obtain the required column permutation.
(no proof)

Lemma : A satisfies CSP $\Rightarrow A \in TU$

Given a profile $P = (\succ_i)_{i \in N}$, consider the matrix $A \in \{0,1\}^{m \times m}$ as constructed as follows :
- Columns $\leftrightarrow \forall\, a \in A$
- rows $\leftrightarrow_{x,k}$ $\forall\, x \in V$, $k \in [m]$
The entry $A(v_{x,k}, C_a) = 1$ if $\text{rank}_x(a) \leq k$.

Theorem Given a preference profile $P = (\succ_i)_{i \in N}$, P is single-peaked iff the matrix A Constructed for P, for the required permutation post the columns so that the 1's in each row are consecutive, given the required base order $\triangleright$.

Proof : Suppose P is single-peaked for some base order $\triangleright$.
Order the columns (alternatives) according to $\triangleright$. Consider any row $v_{x,k}$. Then row has 1's for exactly the top k alternatives of voter x. By the above, these form an interval in $\triangleright$, hence the 1's are contiguous.
Now suppose A satisfies CSP, & let $\triangleright$ be the ordering of these in each row, the 1's are consecutive. Then for each voter x, $k \in [m]$, the top k candidates form an interval in $\triangleright$. By above, P is single-peaked wrt $\triangleright$, & thus holds for all $\succ$. ∎

So why are single-peaked profiles useful?
(Black 1948)
Theorem Let P be single-peaked wrt $\triangleright$, & voters n odd. Then:
(1) If a beat $\succeq_{\triangleright}$ & beats c, then a beats c

(2) Let $p_1, \dots, p_n$ be the top-ranked candidates for the voters, & let p_{med} be the median candidate according to $\triangleright$. Then p_{med} is the Condorcet winner.

For each voter i, let p_i (the peak of i) be i's top-ranked alternative.

Proof : Will show first that p_{med} is the Condorcet winner.
Note that since n is odd, at least $\lceil n/2 \rceil$ voters have top-ranked alternative (peaks) p_i s.t. $p_i \trianglerighteq p_{med}$. & similarly at least $\lceil n/2 \rceil$ voters have peaks p_i s.t. $p_{med} \trianglerighteq p_i$.
Now consider any $a \neq p_{med}$, say $a \triangleright p_{med}$. But then for at least $\frac{n}{2}$ voter p_i, $a \triangleright p_{med} \trianglerighteq p_i$, hence $p_{med} \succ_i a$.

p_{med} beats a. Have p_{med} beats every alternative, & is a Condorcet winner.

(1) Fix any a, b, c, s.t. a beats b, b beats c. Basis: P to $\{a,b,c\}$, then the restricted profile is still single-peaked wrt $\triangleright$ restricted to $\{a,b,c\}$ (Obvious) by Thm.
By (2), there is a Condorcet winner in $P_{|abc}$. Thus has to be a, since a beats b, b beats c. But then a beats c. ∎

Thus, consider the directed graph $G = (N, E)$ where $\forall\, a, b \in A$, there is an edge (a,b) if a beats b, & (b,a) if b beats a (exactly one of these edges must exist). Call this graph the majority graph, for now that
(1) this is a tournament: there is a directed edge for every pair of vertices, and
(2) by the theorem, the graph is acyclic, & it has a Hamiltonian path.

For alternatives A, fix a base order $\triangleright$. Let $SP(\triangleright)$ be the set of all single-peaked rankings wrt to $\triangleright$. Suppose all these preferences of the voters are in $SP(\triangleright)$.
Consider the SCF $f_{med} : SP(\triangleright)^n \to A$, defined as follows:

Median Rule: Given a preference profile P, single-peaked wrt $\triangleright$, let $p_1, \dots, p_n$ be the peaks of the voters, and order the voters so that $p_1 \trianglerighteq p_2 \trianglerighteq \dots \trianglerighteq p_n$. Then: if n is odd, $f_{med} = p_{\frac{n+1}{2}}$, else $f_{med} = p_{\frac{n}{2}}$. f_{med} returns f_{med}.

Theorem If rule is repeated preferences are in $SP(\triangleright)$, then f_{med} is both non-dictatorial & strategyproof.

Proof: **non-dictatorial**: If not voter report a non-peak, then f_{med} returns a non-dictatorial. Suppose $a \triangleright \dots \triangleright b$. Then for voter i let a be peak for voter i, & b be the peak for all the other voters. Then f_{med} returns a.

strategyproof : Let p_i be i's true peak & P be an preference profile where i reports a truthfully & P_{med} be the output. If $p_{med} = p_i$, then i already got the best outcome. Now suppose $p_{med} \triangleright p_i$.
For any report a s.t. $p_{med} \triangleright a$, the number of the peaks must not change. If for any report a s.t. $a \triangleright p_{med}$ the new median $p'_{med} \trianglerighteq p_{med}$ but this is worse for i. Hence misreporting cannot improve the outcome for i. ∎

Generalized medians
Fix alternatives A, base order $\triangleright$. For any alternatives $x = (x_1, \dots, x_m)$ from A. (repetitions allowed).
The generalized median voting rule is
$f_M(P) = \text{median}(p_1, \dots, p_n, x_1, \dots, x_m)$
Let k, r be the first & last alternatives in $\triangleright$.
Think: $\forall\, x_1 \dots x_m$ $k \triangleright x_1 \triangleright \dots \triangleright x_m \triangleright r$

$f_M(P)$ returned k / r

n odd:
$i \in \{1, \dots, \frac{m}{2}\}$: $x_i = k$, $i > \frac{m}{2}$ $x_i = r$
$f_M(P)$ returns median of P

n even:
$i \in \{1, \dots, \frac{m}{2}\}$: $x_i = k$, $i > \frac{m}{2}+1$, $x_i = r$

$f_M(P)$ returns the left median of P.

Theorem Every generalized median rule is strategyproof on $SP(\triangleright)^n$